

# Deterministic Finite Automata

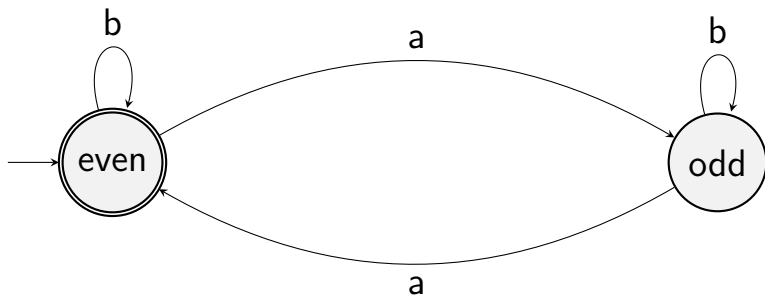
Arjun Chandrasekhar

# Deterministic Finite Automata

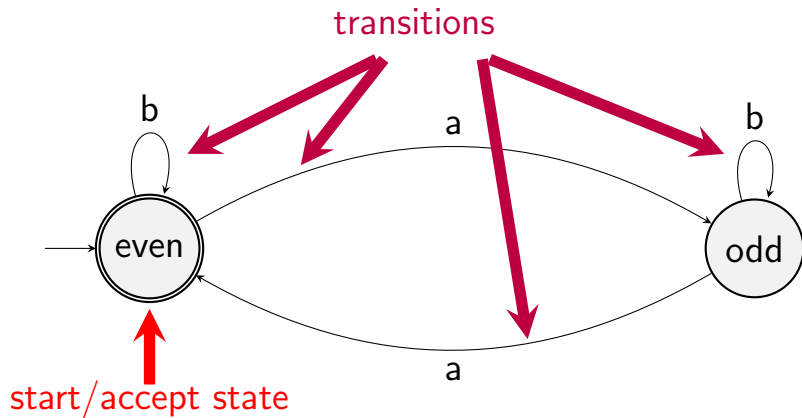
A model of computation that uses a constant amount of memory to recognize a language.

- ▶ Several states
  - ▶ One *start state*
  - ▶ Some states are labelled as *accept states*
- ▶ Transitions (arrows) between states
  - ▶ Each transition is labelled by a character from the alphabet

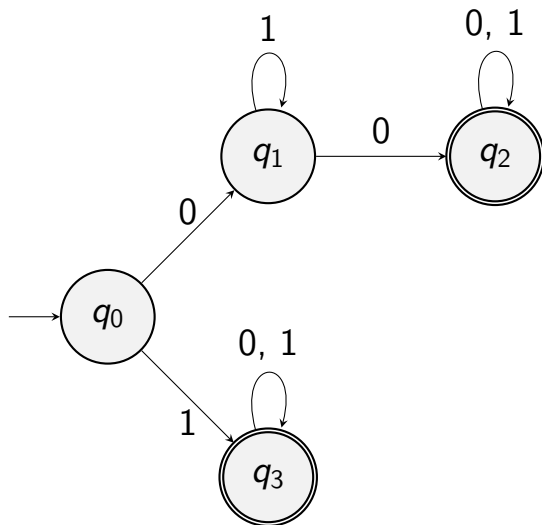
# Deterministic Finite Automata



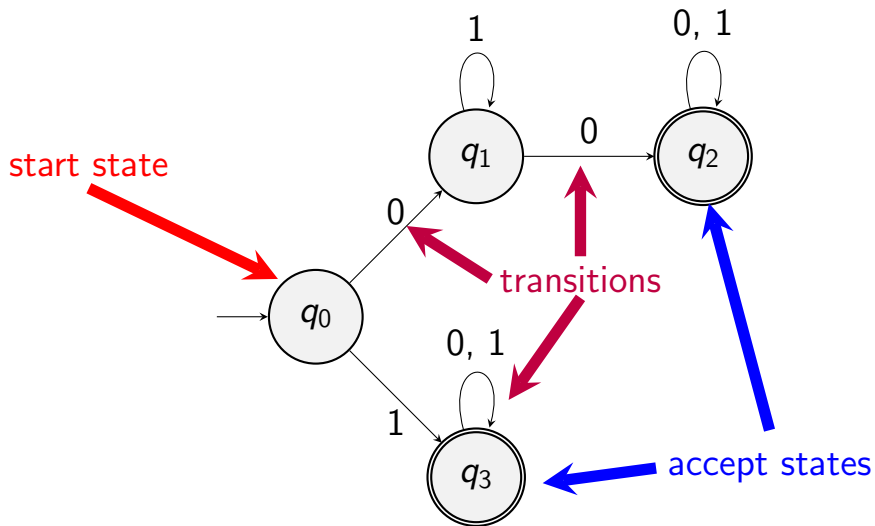
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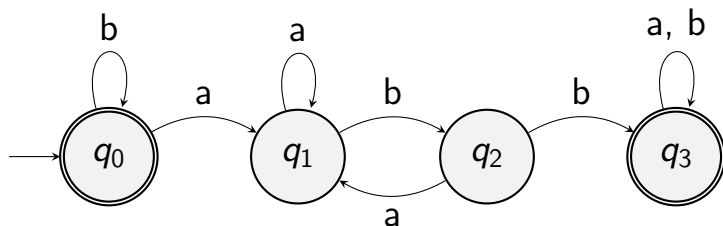
# Deterministic Finite Automata



# Deterministic Finite Automata



# Deterministic finite automata



What is the start state of this DFA?

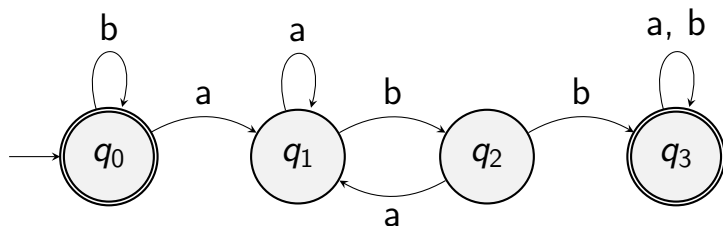
**A.**  $q_0$

**C.**  $q_2$

**B.**  $q_1$

**D.**  $q_3$

# Deterministic finite automata



What is the start state of this DFA?

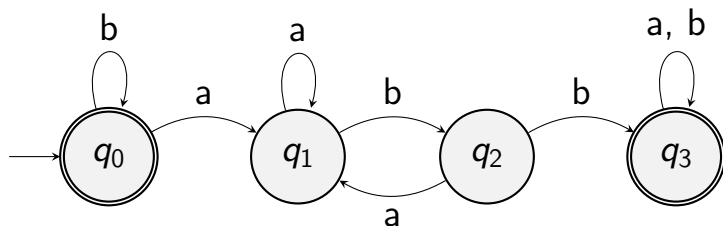
**A.**  $q_0$  ✓

**C.**  $q_2$

**B.**  $q_1$

**D.**  $q_3$

# Deterministic finite automata



What are the accept states of this DFA?

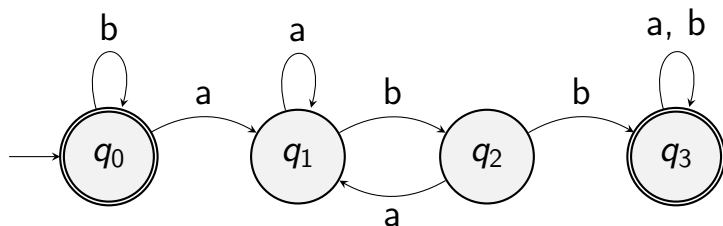
**A.**  $q_0$

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# Deterministic finite automata



What are the accept states of this DFA?

**A.**  $q_0$  ✓

**C.**  $q_2$

**B.**  $q_1$

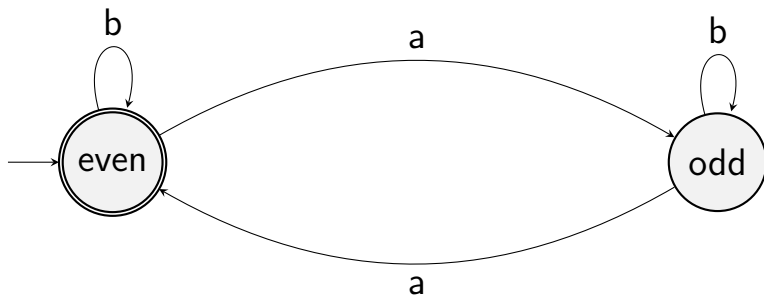
**D.**  $q_3$  ✓

# Computation on a DFA

- ▶ Start in the start state
- ▶ Scan the symbols in the input one by one
- ▶ For each symbol  $\sigma$  scanned:
  - ▶ Go to the next state by following the arrow with the label  $\sigma$
- ▶ After scanning all of the input, if the DFA is in an accept state, the input is accepted. Otherwise the input is rejected

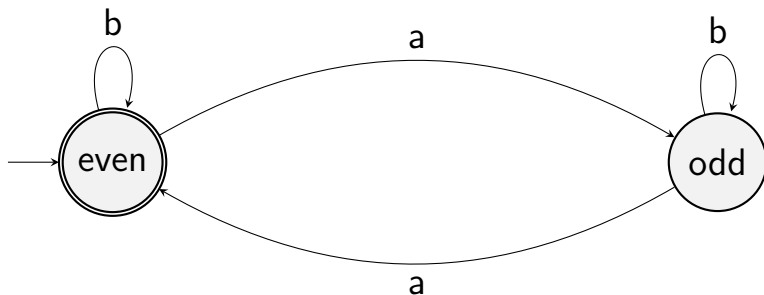
# Deterministic finite automata

What does this DFA do on input aaaa?



# Deterministic finite automata

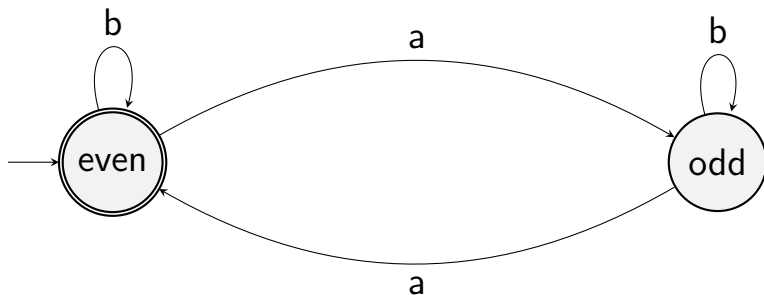
What does this DFA do on input aaaa?



even  $\xrightarrow{a}$  odd  $\xrightarrow{a}$  even  $\xrightarrow{a}$  odd  $\xrightarrow{a}$  even  $\rightarrow$  ACCEPT

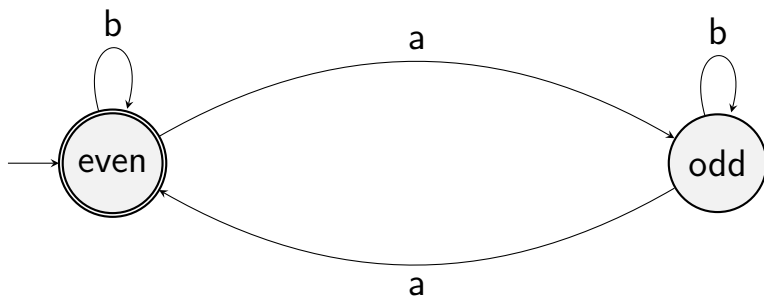
# Deterministic finite automata

What does this DFA do on input b?



# Deterministic finite automata

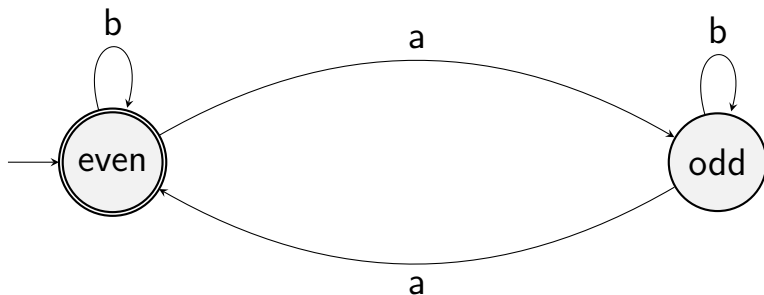
What does this DFA do on input b?



even  $\xrightarrow{b}$  even  $\rightarrow$  ACCEPT

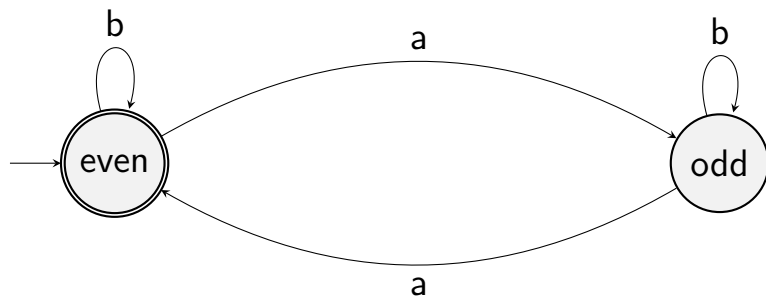
# Deterministic finite automata

What does this DFA do on input abb?



# Deterministic finite automata

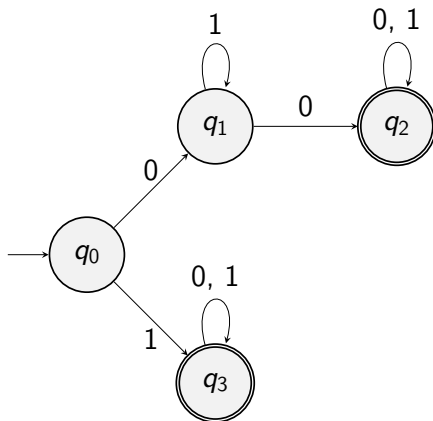
What does this DFA do on input abb?



even  $\xrightarrow{a}$  odd  $\xrightarrow{b}$  odd  $\xrightarrow{b}$  odd  $\rightarrow$  REJECT

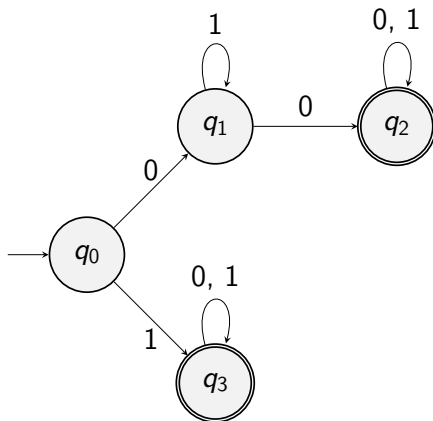
# Deterministic Finite Automata

What does this DFA do on input 0110?



# Deterministic Finite Automata

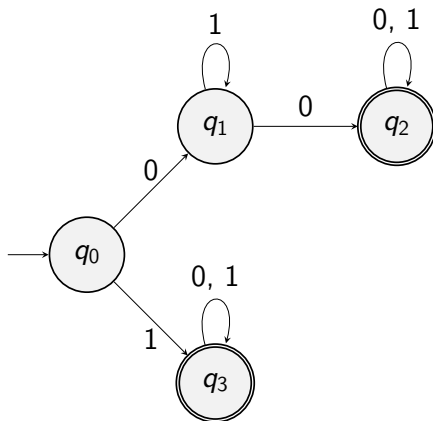
What does this DFA do on input 0110?



$q_0 \xrightarrow{0} q_1 \xrightarrow{1} q_1 \xrightarrow{1} q_1 \xrightarrow{0} q_2 \rightarrow \text{ACCEPT}$  11 / 27

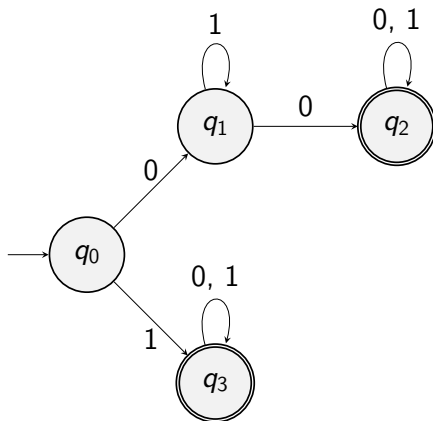
# Deterministic Finite Automata

What does this DFA do on input 1011?



# Deterministic Finite Automata

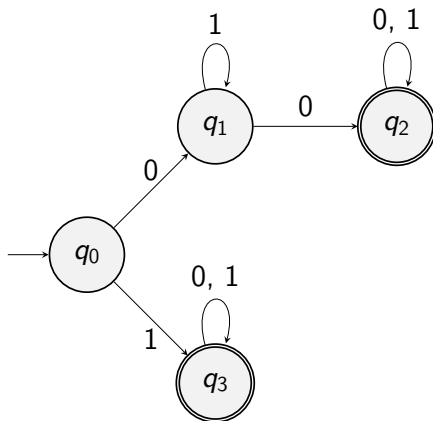
What does this DFA do on input 1011?



$q_0 \xrightarrow{1} q_3 \xrightarrow{0} q_3 \xrightarrow{1} q_3 \xrightarrow{1} q_3 \rightarrow \text{ACCEPT}$

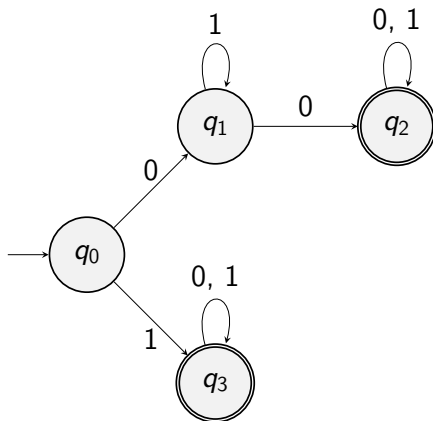
# Deterministic Finite Automata

What does this DFA do on input 0111?



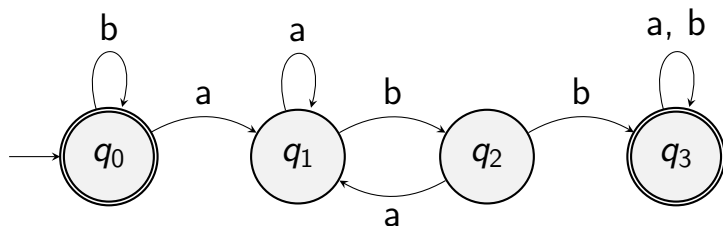
# Deterministic Finite Automata

What does this DFA do on input 0111?



$q_0 \xrightarrow{0} q_1 \xrightarrow{1} q_1 \xrightarrow{1} q_1 \xrightarrow{1} q_1 \rightarrow \text{REJECT}$  13 / 27

# Deterministic Finite Automata



What strings are accepted by this DFA?

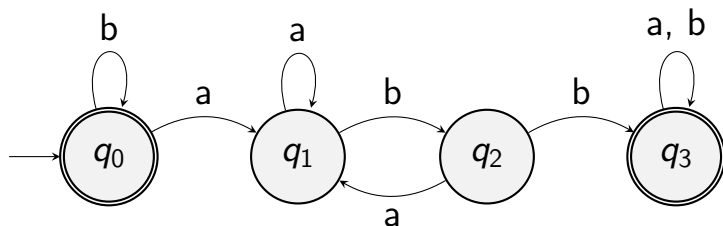
**A.**  $\epsilon$

**C.**  $bbbb a$

**B.**  $abbbb$

**D.**  $b^{100} = \underbrace{b \dots b}_{100}$

# Deterministic Finite Automata



What strings are accepted by this DFA?

**A.**  $\epsilon$  ✓

**C.**  $bbbb a$

**B.**  $abbbb$  ✓

**D.**  $b^{100} = \underbrace{b \dots b}_{100}$  ✓

# Formal definition of a DFA

**Def:** A **Deterministic Finite Automata (DFA)** is a 5-tuple  $(Q, \Sigma, \delta, q_s, F)$

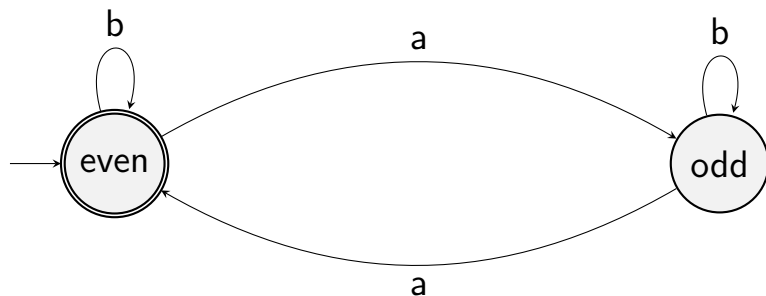
- ▶  $Q$  - The set of states in the DFA
- ▶  $\Sigma$  - The alphabet of characters that the DFA can read
- ▶  $\delta : Q \times \Sigma \rightarrow Q$  - the transition function
  - ▶ The rules for how to move between states
  - ▶ **Input:** Current state & next symbol
  - ▶ **Output:** Next state
- ▶  $q_s$  - the starting state
- ▶  $F$  - the set of accepting states

# Formal definition DFA computation

We say a DFA  $D = (Q, \Sigma, q_s, \delta, F)$  accepts a string  $w = \sigma_1\sigma_2 \dots \sigma_n$  if there exists a sequence of states  $q_0, q_1, \dots, q_n$  such that:

- ▶  $q_0 = q_s$  (start in the start state)
- ▶  $q_i = \delta(q_{i-1}, \sigma_i)$  for all  $i$  (all transitions are valid)
- ▶  $q_n \in F$  (finish in an accept state)

# Formal definition example



►  $Q = \{\text{even}, \text{odd}\}$

►  $\Sigma = \{a, b\}$

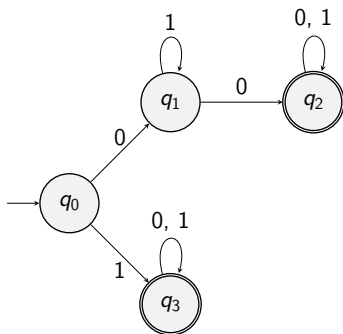
►  $q_s = \text{even}$

►  $F = \{\text{even}\}$

►  $\delta =$

|      | a    | b    |
|------|------|------|
| even | odd  | even |
| odd  | even | odd  |

# Write the formal definition



►  $Q = \{q_0, q_1, q_2, q_3\}$

►  $\Sigma = \{0, 1\}$

►  $q_s = q_0$

►  $F = \{q_2, q_3\}$

►  $\delta =$

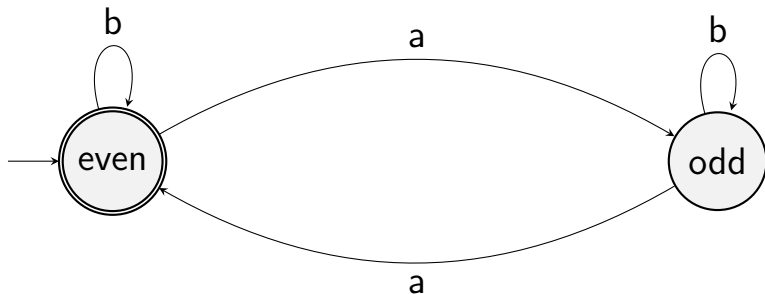
|       | 0     | 1     |
|-------|-------|-------|
| $q_0$ | $q_1$ | $q_3$ |
| $q_1$ | $q_2$ | $q_1$ |
| $q_2$ | $q_2$ | $q_2$ |
| $q_3$ | $q_3$ | $q_3$ |

# The language of a DFA

- ▶ Let  $D$  be a DFA
- ▶ We say a DFA  $D$  **recognizes** a language  $L$  if for all  $w \in \Sigma^*$ :
  - ▶ If  $w \in L$ , then  $D$  accepts  $w$
  - ▶ If  $w \notin L$ , then  $D$  rejects  $w$
  - ▶ Do we ever have to worry about  $D$  looping?
- ▶ The **language of  $D$** , denoted  $L(D)$  is the (unique) language that  $D$  recognizes - that is, the set of all strings that  $D$  accepts

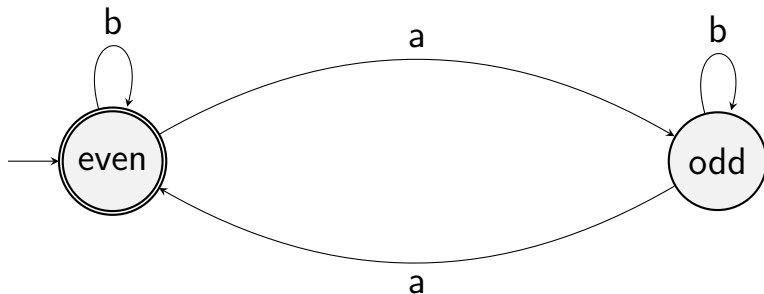
# The language of a DFA

What the language of this DFA?



# The language of a DFA

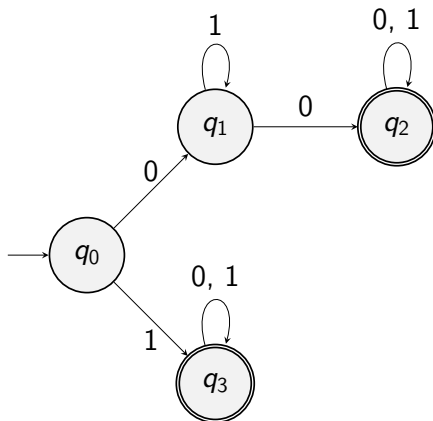
What the language of this DFA?



$$L(D) = \{w \mid w \text{ has an even number of } a\text{'s}\}$$

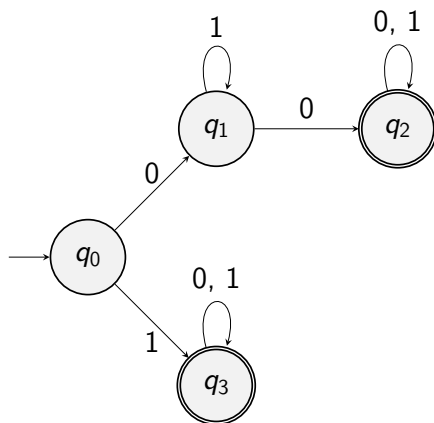
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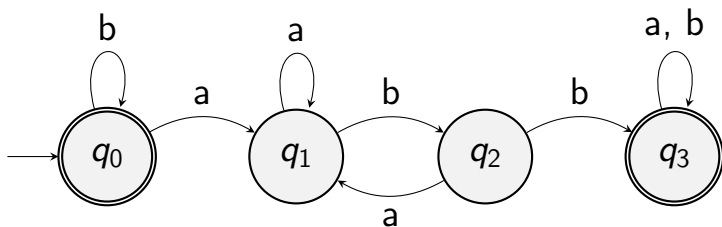
What the language of this DFA?



$$L(D) = \{w \mid w \text{ starts with } 1 \text{ or has at least two } 0\text{'s}\}$$

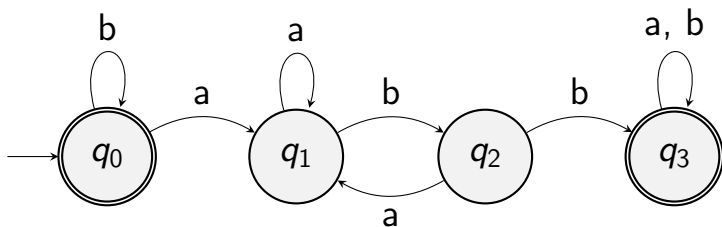
# The language of a DFA

What the language of this DFA?



# The language of a DFA

What the language of this DFA?



$$L(Q) = \{w \mid w \text{ contains only b's or contains abb}\}$$

# Deterministic Finite Automata

Some comments:

- ▶ A transition must be defined for every symbol in the alphabet, for every state in the DFA.
- ▶ There can be more than one accepting state.
- ▶ The set of languages that can be recognized by some DFA is called the **regular languages**.
  - ▶ A language  $L$  is regular if and only if some DFA  $D$  recognizes  $L$

# Design a DFA

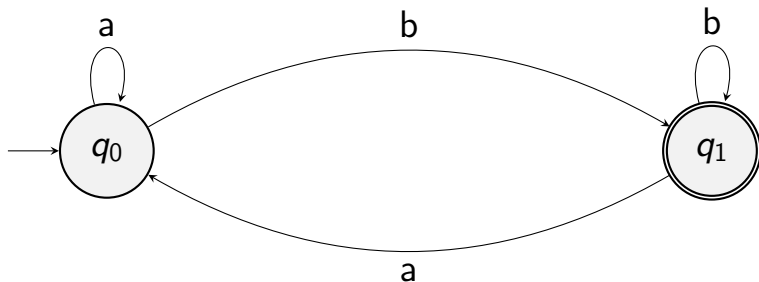
Let  $\Sigma = \{a, b\}$ . Let's show that the following language is regular:

$$L = \{w \mid w \text{ ends with } b\}$$

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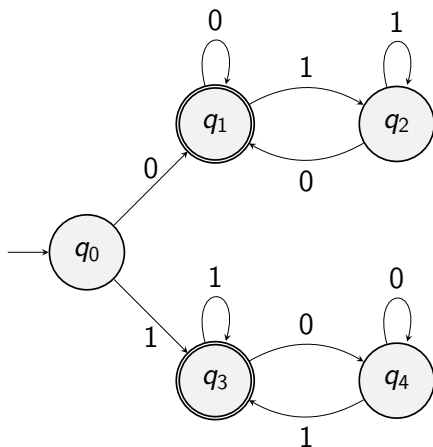
Let  $\Sigma = \{0, 1\}$ . Let's show that the following language is regular:

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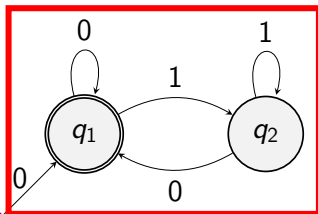


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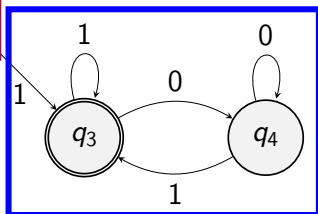
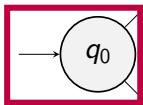
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should we  
accept  $\epsilon$ ?



starts/ends  
with 0



starts/ends  
with 1

# Design a DFA

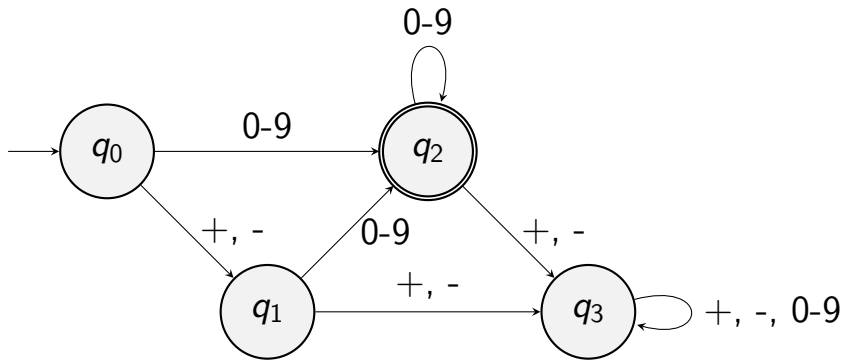
Let  $\Sigma = \{+, -, 0, 1, \dots, 9\}$ . Let's show that the following language is regular:

$$L = \{w \mid w \text{ is a valid integer literal}\}$$

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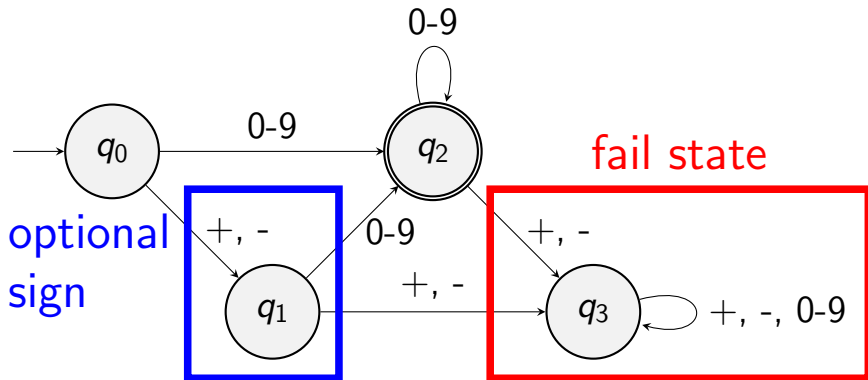
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$$L = \{w \mid w \text{ is a valid integer literal}\}$$



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Let  $\Sigma = \{a, b\}$ . Let's show that the following language is regular:

$$L = \{w \mid w \text{ ends in } bba\}$$

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